# Artificial Intelligence ENCS 434

# **Adversarial Search & Games**



# Game Playing and AI

Why would game playing be a good problem for AI research?

- $\bullet$  Game-playing is non-trivial
	- Need to display "human-like" intelligence
	- Some games (such as chess) are very complex
	- Requires decision-making within a time-limit
		- More realistic than other search problems
- Games are played in a controlled environment
	- Can do experiments, repeat games, etc
	- Good for evaluating research systems
- Can compare humans and computers directly
	- Can evaluate percentage of wins/losses to quantify performance
- $\triangleright$  All the information is available
	- Human and computer have equal information

# How Does a Computer Play a Game?

- $\Box$  A way to play a game is to:
	- Consider all the legal moves you can make
	- Compute the new position resulting from each move
	- Evaluate each resulting position and determine which is best
	- $\bullet$  Make that move
	- Wait for your opponent to move and repeat
- $\Box$  Key problems are:
	- $\bullet$  Representing the "board"
	- Generating all next legal boards
	- $\supset$  Evaluating a position

# Tic-Tac-Toe Game

- Tic-Tac-Toe
	- $\bullet$  b ~ 5 legal moves,
	- $\bullet$  d ~ total of 9 moves
	- $\bullet$  5<sup>9</sup> = 1,953,125
	- $\bullet$  9! = 362,880 (Computer goes first)
	- $\bullet$  8! = 40,320 (Computer goes second)

### **Game Playing: Adversarial Search**

#### **Introduction**

**Different kinds of games:** 



- Games with perfect information. No randomness is involved.
- Games with imperfect information. Random factors are part of the game.

### Games as Adversarial Search

- many games can be formulated as search problems
- the zero-sum utility function leads to an adversarial situation
	- in order for one agent to win, the other necessarily has to lose
- factors complicating the search task
	- potentially huge search spaces
	- elements of chance
	- multi-person games, teams
	- time limits
	- imprecise rules

# Difficulties with Games

- games can be very hard search problems
	- yet reasonably easy to formalize
	- finding the *optimal* solution may be impractical
		- a solution that beats the opponent is "good enough"
	- unforgiving
		- a solution that is "not good enough" not only leads to higher costs, but to a loss to the opponent
- example: chess
	- size of the search space
		- branching factor around 35
		- about 50 moves per player
		- about  $35^{100}$  or  $10^{154}$  nodes
			- about  $10^{40}$  *distinct* nodes (size of the search graph)

# Single-Person Game

- conventional search problem
	- identify a sequence of moves that leads to a winning state
	- examples: Solitaire, dragons and dungeons, Rubik's cube
	- little attention in AI
- some games can be quite challenging
	- some versions of solitaire
	- a heuristic for Rubik's cube was found by the Absolver program

# Searching in a two player game

- Traditional (single agent) search methods only consider how close the agent is to the goal state (e.g. best first search).
- In two player games, decisions of both agents have to be taken into account: a decision made by one agent will affect the resulting search space that the other agent would need to explore.
- Question: Do we have randomness here since the decision made by the opponent is NOT known in advance?
- $\bullet$   $\odot$  No. Not if *all* the moves or choices that the opponent can make are finite and can be known in advance.

# Searching in a two player game

To formalize a two player game as a search problem an agent can be called **MAX** and the opponent can be called **MIN.**

#### **Problem Formulation:**

- **Initial state:** board configurations and the player to move.
- **Successor function:** list of pairs (move, state) specifying legal moves and their resulting states. (moves  $+$  initial state  $=$  game tree)
- **A terminal test:** decide if the game has finished.
- **A utility function:** produces a numerical value for (only) the terminal states. Example: In chess, outcome  $=$  win/loss/draw, with values  $+1$ ,  $-1$ , 0 respectively.
- Players need search tree to determine next move.



• **Idea**: MAX chooses the board with the max utility value, MIN the minimum.

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# MiniMax Algorithm

- Create start node as a MAX node with current board configuration
- Expand nodes down to some depth of lookahead in the game
- Apply the evaluation function at each of the leaf nodes
- $\Box$  "Back up" values for each of the non-leaf nodes until a value is computed for the root node.
	- At MIN nodes, the backed-up value is the minimum of the values associated with its children.
	- At MAX nodes, the backed-up value is the maximum of the values associated with its children.
- Pick the operator associated with the child node whose backed-up value ⊔ determined the value at the root.



terminal nodes: values calculated from the utility function



other nodes: values calculated via minimax algorithm









# MiniMax Properties

Assume all terminal states are at depth d

<sup>S</sup> Space complexity?

Depth-first search, so  $O(bd)$ 

<sup>T</sup>ime complexity?

Given branching factor b, so  $O(b^d)$ 

Time complexity is a major problem!  $\ast$ 

Computer typically only has a finite amount of time to make a move.

- Direct mini-max also is impractical in practice
- Static Board Evaluator (SBE) function \*

Uses heuristics to estimate the value of non-terminal states.

# Pruning

#### Discards parts of the search tree

- Guaranteed not to contain good moves
- Guarantee that the solution is not in that branch or sub-tree
	- If both players make optimal decisions, they will never end up in that part of the search tree
- Use pruning to ignore those branches. ❏

#### Certain moves are not considered

- Won't result in a better evaluation value than a move further up in the tree
- They would lead to a less desirable outcome
- $\Box$  Applies to moves by both players
	- $\infty$  (alpha) indicates the best choice for Max so far never decreases
		- Highest Evaluation value seen so far (initialize to -infinity)
	- $\supset \beta$  (beta) indicates the best choice for Min so far never increases
		- Lowest Evaluation value seen so far (initialize to  $+i$ nfinity)

# Alpha-Beta Pruning

Beta cutoff pruning occurs when maximizing if child's alpha  $\geq$  parent's beta Stop expanding children. Why?

Opponent won't allow computer to take this move

 $\Box$  Alpha cutoff pruning occurs when minimizing if parent's alpha  $\geq$  child's beta Stop expanding children. Why?

 $\supset$  Computer has a better move than this



#### $\alpha$  best choice for Max  $\beta$  best choice for Min

- we assume a depth-first, left-to-right search as basic strategy
- the range of the possible values for each node are indicated
	- initially  $[-\infty, +\infty]$
	- from Max's or Min's perspective
	- these *local* values reflect the values of the sub-trees in that node; the *global* values  $\alpha$  and  $\beta$  are the best overall choices so far for Max or Min







- Min obtains the third value from a successor node
- this is the last value from this sub-tree, and the exact value is known
- Max now has a value for its first successor node, but hopes that something better might still come



- Min continues with the next sub-tree, and gets a better value
- Max has a better choice from its perspective, however, and will not consider a move in the sub-tree currently explored by Min
	- $\bullet$  initially  $[-\infty, +\infty]$



this is a case of *pruning*, indicated by





- Min explores the next sub-tree, and finds a value that is worse than the other nodes at this level
- if Min is not able to find something lower, then Max will choose this branch, so Min must explore more successor nodes



- Min is lucky, and finds a value that is the same as the current worst value at this level
- Max can choose this branch, or the other branch with the same value



- Min could continue searching this sub-tree to see if there is a value that is less than the current worst alternative in order to give Max as few choices as possible
	- this depends on the specific implementation
- Max knows the best value for its sub-tree

# Properties of Alpha-Beta Pruning

- in the ideal case, the best successor node is examined first
	- results in  $O(b^{d/2})$  nodes to be searched instead of  $O(b^d)$
	- alpha-beta can look ahead twice as far as minimax
	- in practice, simple ordering functions are quite useful
- assumes an idealized tree model
	- uniform branching factor, path length
	- random distribution of leaf evaluation values
- transpositions tables can be used to store permutations
	- sequences of moves that lead to the same position
- requires additional information for good players
	- game-specific background knowledge
	- empirical data

# Imperfect Decisions

- complete search is impractical for most games
- alternative: search the tree only to a certain depth
	- requires a cutoff-test to determine where to stop
		- replaces the terminal test
		- the nodes at that level effectively become terminal leave nodes
	- uses a heuristics-based evaluation function to estimate the expected utility of the game from those leave nodes

# Evaluation Function

- determines the performance of a game-playing program
- must be consistent with the utility function
	- values for terminal nodes (or at least their order) must be the same
- tradeoff between accuracy and time cost
	- without time limits, minimax could be used
- should reflect the actual chances of winning
- frequently weighted linear functions are used
	- $E = w_l f_l + w_2 f_2 + \ldots + w_n f_n$
	- $\bullet$  combination of features, weighted by their relevance

### Example: Tic-Tac-Toe

simple evaluation function

 $E(s) = (rx + cx + dx) - (ro + co + do)$ 

(number of rows, columns, and diagonals open for  $MAX$ ) – (number of rows, columns, and diagonals open for MIN )

- 1-ply lookahead
	- start at the top of the tree
	- evaluate all 9 choices for player 1
	- pick the maximum E-value
- 2-ply lookahead
	- also looks at the opponents possible move
		- assuming that the opponents picks the minimum E-value





### Checkers Case Study

- initial board configuration
	- Black single on 20
		-
	- single on 21 king on 31 • Red single on 23 king on 22
	- evaluation function

$$
E(s) = (5 x_1 + x_2) - (5r_1 + r_2)
$$

where

- $x_1$  = black king advantage,
- $x_2$  = black single advantage,
- $r_1$  = red king advantage,
- $r_2$  = red single advantage























# Search Limits

- search must be cut off because of time or space limitations
- strategies like depth-limited or iterative deepening search can be used
	- don't take advantage of knowledge about the problem
- more refined strategies apply background knowledge
	- quiescent search
		- cut off only parts of the search space that don't exhibit big changes in the evaluation function

# Horizon Problem

- moves may have disastrous consequences in the future, but the consequences are not visible
	- the corresponding change in the evaluation function will only become evident at deeper levels
		- they are "beyond the horizon"
- determining the horizon is an open problem without a general solution
	- only some pragmatic approaches restricted to specific games or situation

### Games with Chance

- in many games, there is a degree of unpredictability through random elements
	- throwing dice, card distribution, roulette wheel, ...
- this requires *chance nodes* in addition to the Max and Min nodes
	- branches indicate possible variations
	- each branch indicates the outcome and its likelihood

# Decisions with Chance

- the utility value of a position depends on the random element
	- the definite minimax value must be replaced by an expected value
- calculation of expected values
	- utility function for terminal nodes
	- for all other nodes
		- calculate the utility for each chance event
		- weigh by the chance that the event occurs
		- add up the individual utilities

# Chapter Summary

- many game techniques are derived from search methods
- the minimax algorithm determines the best move for a player by calculating the complete game tree
- alpha-beta pruning dismisses parts of the search tree that are provably irrelevant
- an evaluation function gives an estimate of the utility of a state when a complete search is impractical
- chance events can be incorporated into the minimax algorithm by considering the weighted probabilities of chance events