



## Chapter 6

# Binary Search Trees

1<sup>st</sup> semester

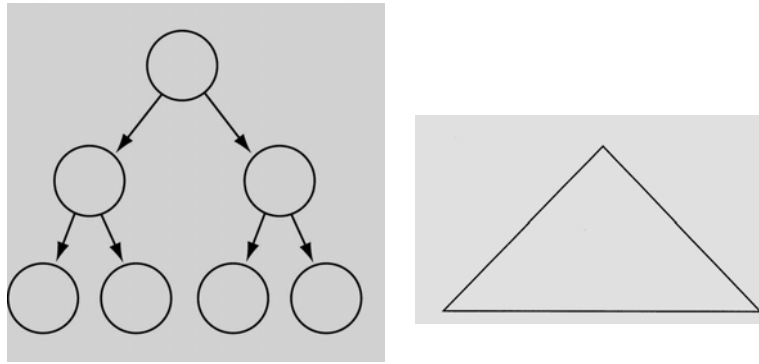
2007/2008

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# Heaps

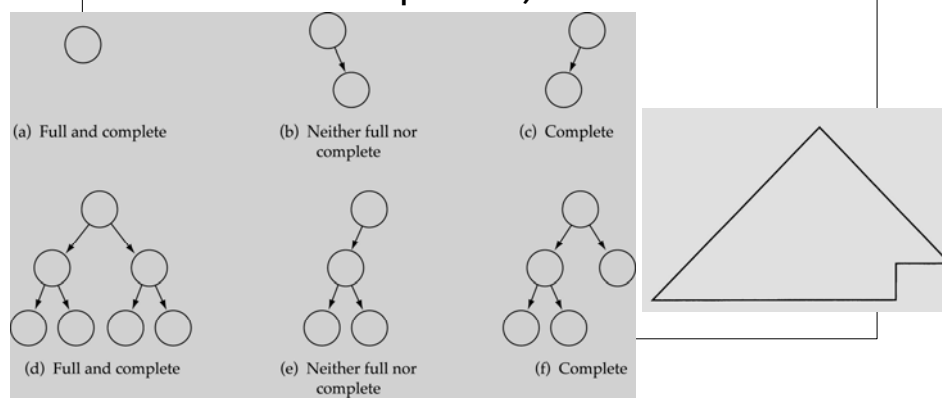
# Full Binary Tree

- Every non-leaf node has two children
- All the leaves are on the same level



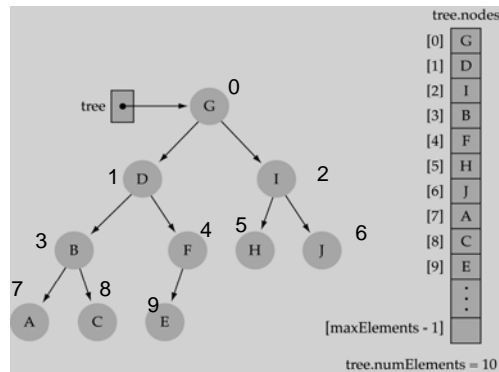
# Complete Binary Tree

- A binary tree that is either full or full through the next-to-last level
- The last level is full from left to right (i.e., leaves are as far to the left as possible)



## Array-based representation of complete binary trees

- Preserve parent-child relationships by storing the tree elements in the array

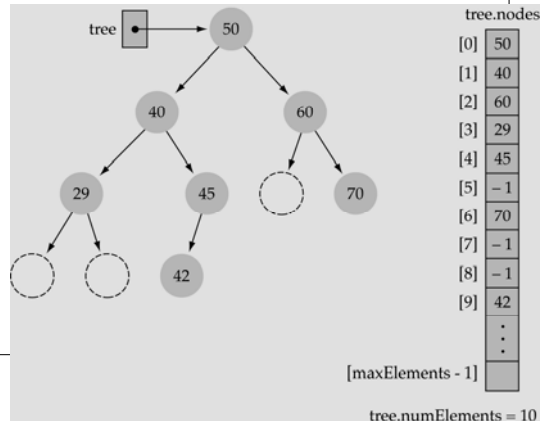


## Array-based representation of complete binary trees

- Parent-child relationships:
  - left child of  $tree.nodes[index] = tree.nodes[2*index+1]$
  - right child of  $tree.nodes[index] = tree.nodes[2*index+2]$
  - parent node of  $tree.nodes[index] = tree.nodes[(index-1)/2]$   
(int division-truncate)
- Leaf nodes:
  - $tree.nodes[numElements/2]$  to  $tree.nodes[numElements - 1]$

## Array-based representation of complete binary trees

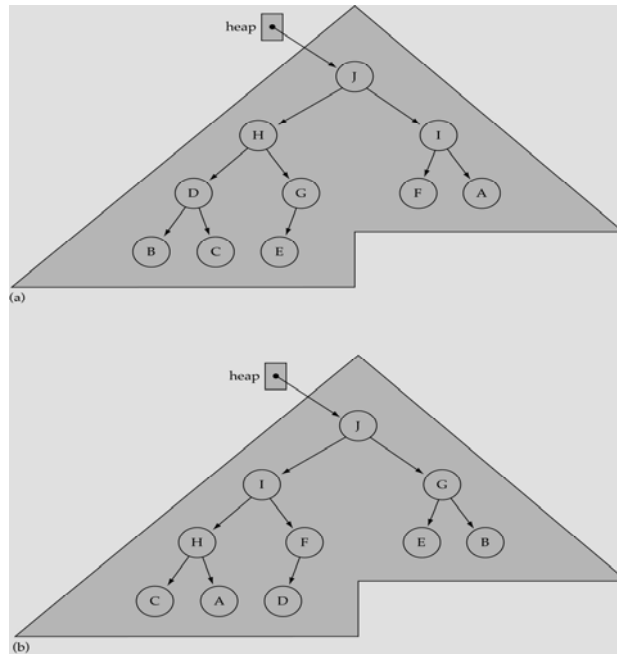
- Full or complete trees can be implemented easily using an array-based representation (elements occupy contiguous array slots)
- "Dummy nodes" are required for trees which are not full or complete



## What is a heap?

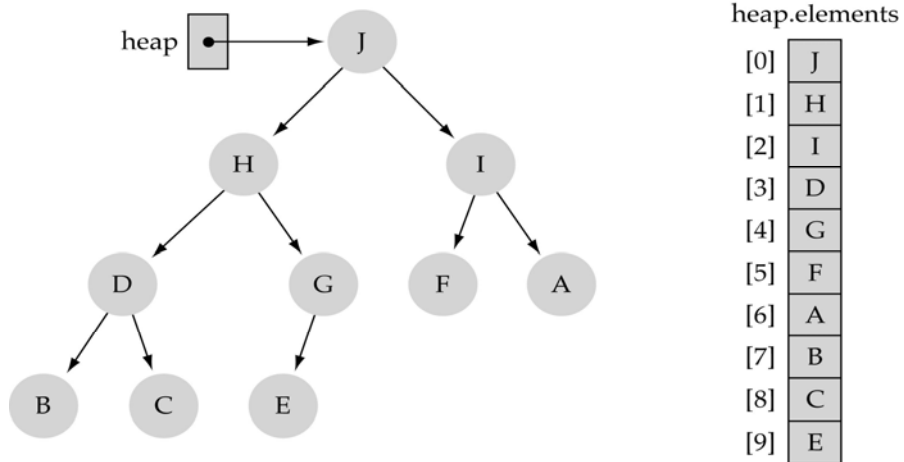
- It is a binary tree with the following properties:
  - *Property 1*: it is a complete binary tree
  - *Property 2*: the value stored at a node is greater or equal to the values stored at the children

# What is a heap? (cont.)



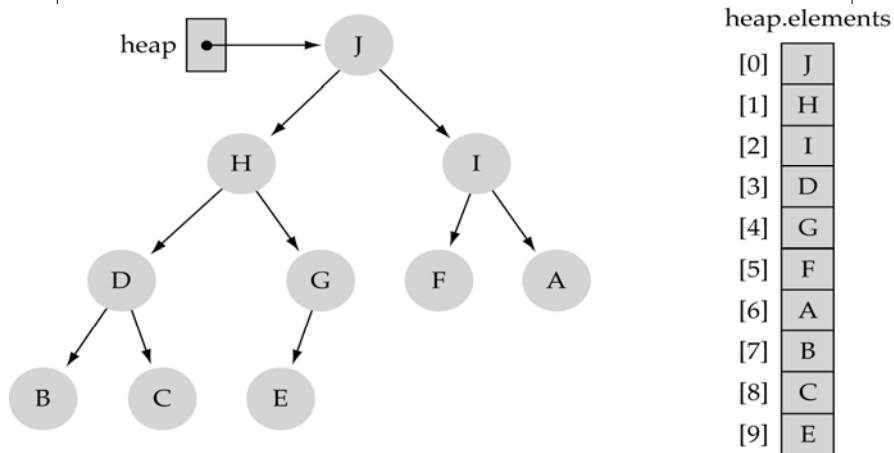
## Largest heap element

- From *Property 2*, the largest value of the heap is always stored at the root



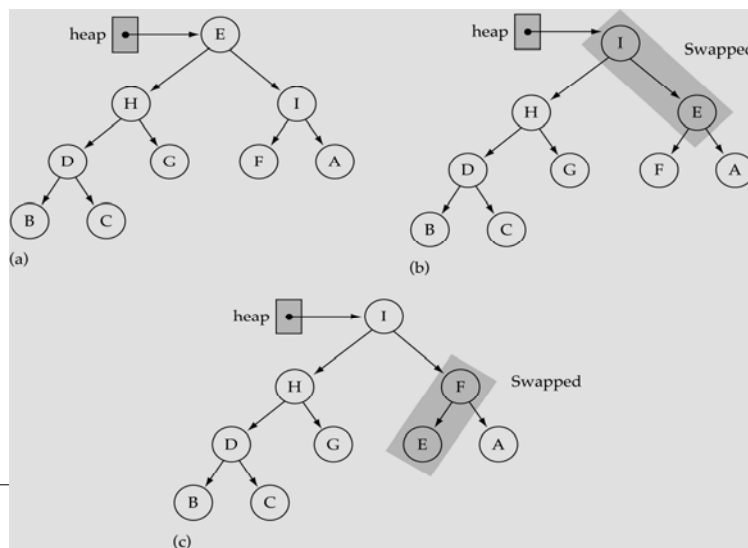
# Heap implementation using array representation

- A heap is a complete binary tree, so it is easy to be implemented using an array representation



# The ReheapDown function (used by deleteltem)

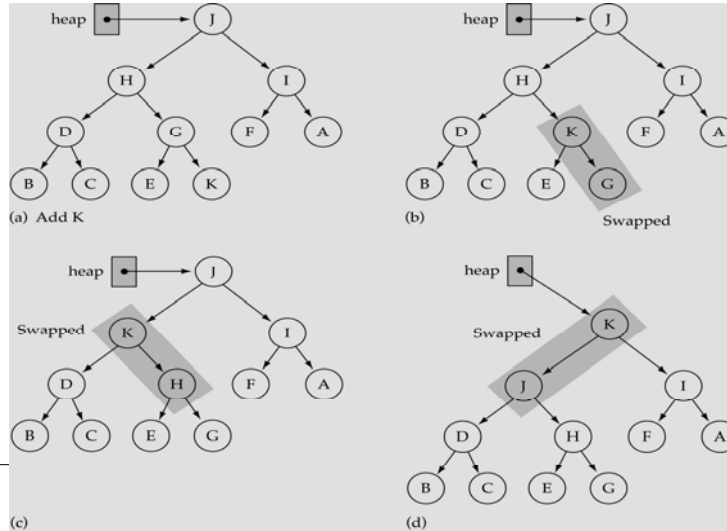
- Assumption: heap property is violated at the root of the tree



# The ReheapUp function

(used by insertItem)

- Assumption: heap property is violated at the rightmost node at the last level of the tree



# ReheapDown function

rightmost node  
in the last level

```


template<class ItemType>
void HeapType<ItemType>::ReheapDown(int root, int bottom) {
    int maxChild, rightChild, leftChild;

    leftChild = 2*root+1;
    rightChild = 2*root+2;

    if(leftChild <= bottom) { // left child is part of the heap
        if(leftChild == bottom) // only one child
            maxChild = leftChild;
        else {
            if(elements[leftChild] <= elements[rightChild])
                maxChild = rightChild;
            else
                maxChild = leftChild;
        }
        if(elements[root] < elements[maxChild]) {
            Swap(elements, root, maxChild);
            ReheapDown(maxChild, bottom);
        }
    }
}
    
```

# ReheapUp function

Assumption:  
heap property  
is violated at bottom



```
template<class ItemType>
void HeapType<ItemType>::ReheapUp(int
    root, int bottom){
    int parent;

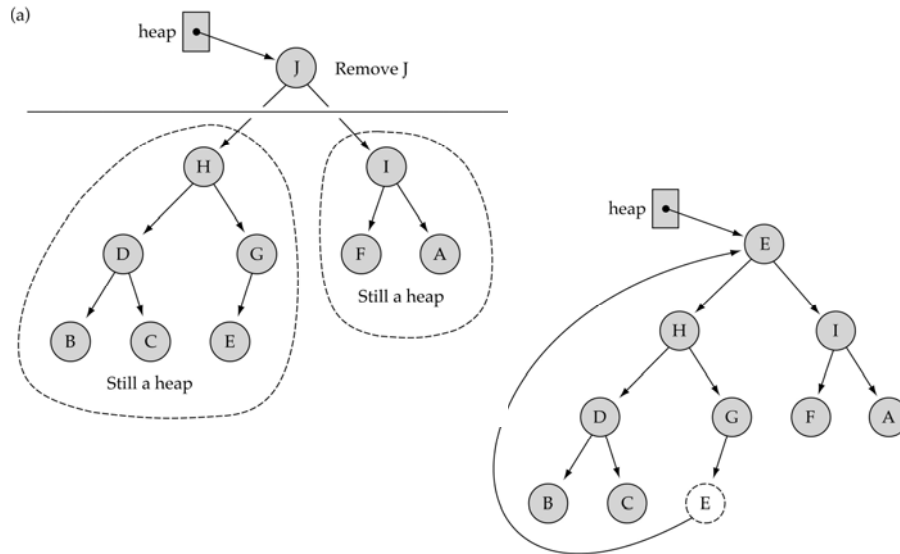
    if(bottom > root) { // tree is not empty
        parent = (bottom-1)/2;
        if(elements[parent] < elements[bottom]) {
            Swap(elements, parent, bottom);
            ReheapUp(root, parent);
        }
    }
}
```

## Removing the largest element from the heap

- 1) Copy the bottom rightmost element to the root
- 2) Delete the bottom rightmost node
- 3) Fix the heap property by calling *ReheapDown*



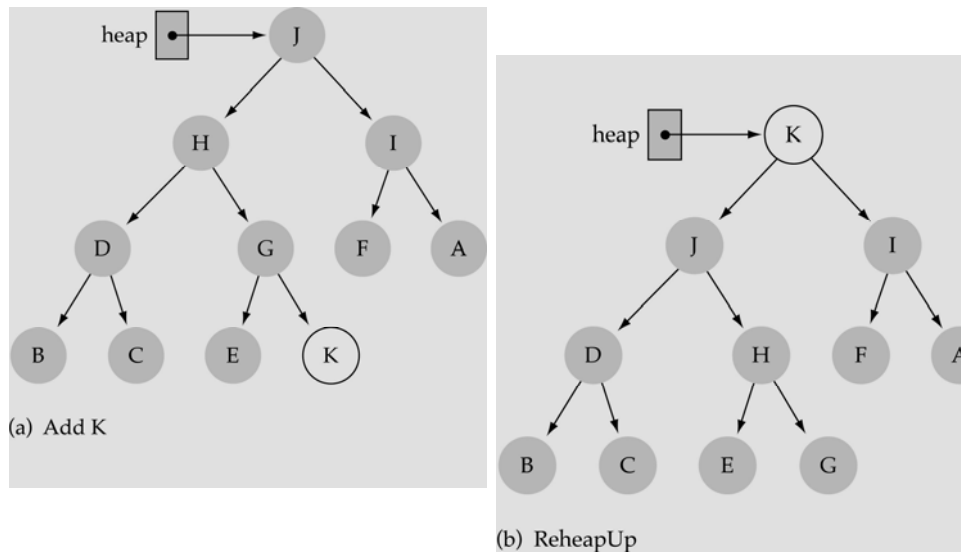
## Removing the largest element from the heap (cont.)



## Inserting a new element into the heap

- 1) Insert the new element in the next bottom leftmost place
- 2) Fix the heap property by calling *ReheapUp*

## Inserting a new element into the heap (cont.)



## Priority Queues

### ■ What is a priority queue?

- It is a queue with each element being associated with a "priority"
- From the elements in the queue, the one with the highest priority is dequeued first