

Relations on Sets

A more formal way to refer to the kind of relation defined on sets is to call it a **binary relation** because it is a subset of a **Cartesian product of two sets**.

At the end of this section we define an *n*-ary relation to be a subset of a Cartesian product of *n* sets, where *n* is any *integer greater than or equal to two*.

Such a relation is the fundamental structure used in<u>relational databases</u>. However, because we focus on binary relations in this text, when we use the term *relation* by itself, we will <u>mean binary relation</u>.



































Example 6 – Directed Graph of a Relation Let $A = \{3, 4, 5, 6, 7, 8\}$ and define a relation R on A as follows: For all $x, y \in A$, x R y = 2 | (x - y).

Draw the directed graph of R.

Solution:

Note that 3 R 3 because 3 - 3 = 0 and 2 | 0 since $0 = 2 \cdot 0$. Thus there is a loop from 3 to itself.

Similarly, there is a loop from 4 to itself, from 5 to itself, and so forth, since the difference of each integer with itself is 0, and $2 \mid 0$.





Example 7 – A Simple Database

The following is a radically simplified version of a database that might be used in a hospital.

Let A_1 be a set of positive integers, A_2 a set of alphabetic character strings, A_3 a set of numeric character strings, and A_4 a set of alphabetic character strings.

Define a quaternary relation R on $A_1 \times A_2 \times A_3 \times A_4$ as follows:

 $(a_1, a_2, a_3, a_4) \in R$ a patient with patient ID number a_1 , named a_2 , was admitted on date a_3 , with primary diagnosis a_4 .



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Reflexivity, Symmetry, and Transitivity

Let $A = \{2, 3, 4, 6, 7, 9\}$ and define a relation R on A as follows: For all $x, y \in A$, $x R y \Leftrightarrow 3 | (x - y).$ Then 2 R 2 because 2 – 2 = 0, and 3 | 0.

Similarly, 3 R 3, 4 R 4, 6 R 6, 7 R 7, and 9 R 9.

Also 6 R 3 because 6 - 3 = 3, and 3 | 3.

And 3 R 6 because 3 - 6 = -(6 - 3) = -3, and 3 | (-3).

Similarly, 3 R 9, 9 R 3, 6 R 9, 9 R 6, 4 R 7, and 7 R 4.

Reflexivity, Symmetry, and Transitivity





Reflexivity, Symmetry, and Transitivity

Because of the equivalence of the expressions x R y and $(x, y) \in R$ for all x and y in A, the reflexive, symmetric, and transitive properties can also be written as follows:

1. <i>R</i> is reflexive	\Leftrightarrow	for all x in A , $(x, x) \in R$.
2. <i>R</i> is symmetric	\Leftrightarrow	for all x and y in A, if $(x, y) \in R$ then $(y, x) \in R$.
3. <i>R</i> is transitive	⇔	for all x, y and z in A, if $(x, y) \in R$ and $(y, z) \in R$ then $(x, z) \in R$.

Reflexivity, Symmetry, and Transitivity In informal terms, properties (1)–(3) say the following: 1. <u>Reflexive</u>: Each element is related to itself 2. <u>Symmetric</u>: If any <u>one element</u> is related to any other element, then the <u>second element</u> is related to <u>the first</u>. 3. <u>Transitive</u>: If any <u>one element</u> is related to a <u>second</u> and that <u>second</u> <u>element</u> is related to <u>a third</u>, then the <u>first element is related to the third</u>. Note that the definitions of reflexivity, symmetry, and transitivity are<u>universal statements</u>.

Reflexivity, Symmetry, and Transitivity

This means that to prove a relation has one of the properties, you<u>use either</u> the method of exhaustion or the method of generalizing from the generic particular.

Now consider what it means for a relation*not* to have one of the properties defined previously. We have known that the negation of a universal statement is existential.

Hence if *R* is a relation on a set *A*, then

1. *R* is **not reflexive** there is an element *x* in *A* such that $x \not R x$ [that is, such that $(x, x) \notin R$].

Reflexivity, Symmetry, and Transitivity

2. <i>R</i> is not symmetric	there are elements x and y in $y \mathcal{R} x$ [that is, such that $(x, y) \in R$ but $(y, x) \notin R$].	A such that x R y but
3. <i>R</i> is not transitive and	there are elements x , y and $y R z$ but $x \mathcal{R} z$ [that is, such $(y, z) \in R$ but $(x, z) \notin R$].	z in A such that $x R ythat (x, y) \in R and$
It follows that you can show finding a counterexample.	that a relation does <i>not</i> have or	e of the properties by









Example 1(b) – Solution

<u>S is transitive</u>: There are three cases for which there is an arrow going from one point of the graph to a second and from the second point to a third

Namely, there are arrows going from 0 to 2 and from 2 to 3; there are arrows going from 0 to 0 and from 0 to 2; and there are arrows going from 0 to 0 and from 0 to 3.

In each case there is an arrow going from the first point to the third.(Note again that the "first," "second," and "third" points need not be distinct.)

This means that whenever $(x, y) \in S$ and $(y, z) \in S$, then $(x, z) \in S$, for all x, $y, z \in \{0, 1, 2, 3\}$, and so S is transitive.

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cont'd



Example 1(c) – Solution

The only way for this **to be false** would be for there to exist elements of *A* that make the hypothesis true and the conclusion false.

That is, there would have to be elements x, y, and z in A such that $(x, y) \in T$ and $(y, z) \in T$ and $(x, z) \notin T$.

In other words, there would have to be two ordered pairs inT that have the potential to "link up" by having the *second* element of one pair be the *first* element of the other pair.

But the only elements in *T* are (0, 1) and (2, 3), and these do not have the potential to link up. <u>Hence the hypothesis is never true</u>. It follows that it is impossible for T not to be transitive, and thus T is transitive.

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Properties of Relations on Infinite Set*A*. To **prove** the relation is reflexive, symmetric, or transitive, first write down what is to be proved. For instance, for **symmetry** you need to prove that $\forall x, y \in A$, if x R y then y R x. Then use the definitions of *A* and *R* to rewrite the statement for the particular case in question. For instance, for the <u>"equality"</u> relation on the set of real numbers, the rewritten statement is

Properties of Relations on Infinite Sets

Sometimes the truth of the rewritten statement will be**immediately obvious** (as it is here).

At other times you will need to prove it using the method of **generalizing from the generic particular**.

We begin with the relation of equality, one of the simplest and yet most important relations.

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The Transitive Closure of a Relation

Generally speaking, a <u>relation fails to be transitive</u> because it fails to contain certain ordered pairs.

For example, if (1, 3) and (3, 4) are in a relation*R*, then the pair (1, 4) *must* be in *R* if *R* is to be transitive.

To obtain a transitive relation from one that is not transitive, it is necessary to add ordered pairs.

Roughly speaking, the relation obtained by adding the least number of ordered pairs to ensure transitivity is called the *transitive closure* of the relation.





















The Relation Induced by a Partition

The fact is that *a relation induced by a partition of a set satisfies all three properties*: reflexivity, symmetry, and transitivity.

Theorem 8.3.1

Let A be a set with a partition and let R be the relation induced by the partition. Then R is reflexive, symmetric, and transitive.





Example 2 – Solution	
R is reflexive: Suppose <i>A</i> is a nonempty subset of {1, 2, 3}. [We must show that <i>A</i> R <i>A</i> .]	
It is true to say that the least element of <i>A</i> equals the least element of <i>A</i> . Thus, by definition of <i>R</i> , <i>A</i> R <i>A</i> .	
R is symmetric: Suppose <i>A</i> and <i>B</i> are nonempty subsets of {1, 2, 3} and <i>A</i> R <i>B</i> . <i>[We must show that B</i> R <i>A.]</i>	
Since <i>A</i> R <i>B</i> , the least element of <i>A</i> equals the least element of <i>B</i> .	
But this implies that the least element of <i>B</i> equals the least element of <i>A</i> , and so, by definition of R , <i>B</i> R <i>A</i> .	
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Example 2 – Solution

<u>R</u> is transitive: Suppose *A*, *B*, and *C* are nonempty subsets of {1, 2, 3}, *A* **R** *B*, and *B* **R** *C*. *[We must show that A* **R** *C*.]

Since $A \ \mathbf{R} B$, the least element of A equals the least element of B and since $B \ \mathbf{R}$, the least element of B equals the least element of C.

Thus the least element of A equals the least element of C, and so, by definition of \mathbf{R} , A \mathbf{R} C.

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Equivalence Classes of an Equivalence Relation

Suppose there is an equivalence relation on a certain set. If *a* is any particular element of the set, then one can ask, <u>What is the subset of all elements that are related to a</u>?" This subset is called the <u>equivalence class</u> of <u>a</u>.

• Definition

Suppose *A* is a set and *R* is an equivalence relation on *A*. For each element *a* in *A*, the **equivalence class of** *a*, denoted [*a*] and called the **class of** *a* for short, is the set of all elements *x* in *A* such that *x* is related to *a* by *R*. In symbols: $[a] = \{x \in A \mid x R a\}$

for all $x \in A$, $x \in [a] \Leftrightarrow x R a$.



First find the equivalence class of every element of A. $\begin{bmatrix} 0 \end{bmatrix} = \{x \in A \mid x \ R \ 0 \} = \{0, 4\} \\ \begin{bmatrix} 1 \end{bmatrix} = \{x \in A \mid x \ R \ 0 \} = \{1, 3\} \\ \begin{bmatrix} 2 \end{bmatrix} = \{x \in A \mid x \ R \ 2 \} = \{2\} \\ \begin{bmatrix} 3 \end{bmatrix} = \{x \in A \mid x \ R \ 3 \} = \{1, 3\} \\ \begin{bmatrix} 4 \end{bmatrix} = \{x \in A \mid x \ R \ 4 \} = \{0, 4\}$ Note that [0] = [4] and [1] = [3]. Thus the *distinct* equivalence classes of the relation are $\{0, 4\}, \{1, 3\}, \text{ and } \{2\}.$

Equivalence Classes of an Equivalence Relation

<u>The first lemma</u> says that if two elements of *A* are related by an equivalence relation *R*, then their equivalence classes are the same.

Lemma 8.3.2

Suppose *A* is a set, *R* is an equivalence relation on *A*, and *a* and *b* are elements of *A*. If *a R b*, then [a] = [b].

This lemma says that if a certain condition is satisfied, then [a] = [b]. Now [a] and [b] are *sets*, and two sets are equal if, and only if, each is a subset of the other.

Equivalence Classes of an Equivalence Relation Hence the proof of the lemma consists of two parts: first, a proof that ঀ) ⊆ [b] and second, a proof that [b] ⊆ [a]. To show each subset relation, it is necessary to show that every element in the left-hand set is an element of the ight-hand set. The second lemma says that any two equivalence classes of an equivalence relation are either mutually disjoint or identical Imma 8.3.3 If A is a set, R is an equivalence relation on A, and a and b are elements of A, then either [a] ∩ [b] = ∅ or [a] = [b].



The statement of Lemma 8.3.3 has the form

if p then (q or r),

where *p* is the statement "*A* is a set, *R* is an equivalence relation on *A*, and *a* and *b* are elements of *A*," *q* is the statement " $[a] \cap [b] = \emptyset$," and *r* is the statement "[a] = [b]."

Theorem 8.3.4 The Partition Induced by an Equivalence Relation

If A is a set and R is an equivalence relation on A, then the distinct equivalence classes of R form a partition of A; that is, the union of the equivalence classes is all of A, and the intersection of any two distinct classes is empty.





Exan	nple 10 – Solution	conť d
Therefore	$= \{x \in \mathbb{Z} \mid x - a = 3k, \text{ for some integer} \}$	<i>k</i> }.
merelore,	$[a] = \{x \in \mathbb{Z} \mid x = 3k + a, \text{ for some integer}\}$	<i>k</i> }.
In particular,	$[0] = \{x \in \mathbb{Z} \mid x = 3k + 0, \text{ for some integer} \}$	<i>k</i> }
	$= \{x \in \mathbb{Z} \mid x = 3k, \text{ for some integer } k\}$	
	$= \{\ldots -9, -6, -3, 0, 3, 6, 9, \ldots\},\$	
	$[1] = \{x \in \mathbb{Z} \mid x = 3k + 1, \text{ for some integer}\}$	<i>k</i> }
	$= \{\ldots - 8, -5, -2, 1, 4, 7, 10, \ldots\},\$	
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Example 12(a) – Solution

Suppose (a, b), (c, d), and (e, f) are particular but arbitrarily chosen elements of A such that (a, b) R (c, d) and (c, d) R (e, f).

[We must show that for all (a, b), (c, d), $(e, f) \in A$, if (a, b) R (c, d) and (c, d) R (e, f), then (a, b) R (e, f).]

[We must show that (a, b) R (e, f).] By definition of R,

(1) ad = bc and (2) cf = de.

Since the second elements of all ordered pairs in *A* are nonzero, $b \neq 0$, $d \neq 0$, and $f \neq 0$.









Solution DEFINITIONS A relation R on a set A is **reflexive** if $(a, a) \in R$ for every element $a \in A$. A relation R on a set A is symmetric if $(b,a) \in R$ whenever $(a,b) \in R$ A relation R on a set A is **transitive** if $(a, b) \in R$ and $(b, c) \in R$ implies $(a,c) \in R$ A relation R is an **equivalence relation** if the relation R is transitive, symmetric and reflexive. The **equivalence class** of a is the set of all elements that are relation to a. Notation: $[a]_R$ $[a] = \{ x \in A | x R a \}$ 97

SOLUTION	
$A = \mathbf{Z} \times \mathbf{Z}$	
$R = \{((a,b), (c,d)) \in A \times A a + d = b + c\}$	
(-) Tf. D isfin-	
(a) 10 proof: R is renexive	
PROOF	
Let $(a,b) \in A = \mathbf{Z} \times \mathbf{Z}$.	
By the commutative property of addition:	
a + b = b + a	
However, this then implies that $(a, b) R(a, b)$ by definition of R .	
Since $(a, b) R (a, b)$ for all $(a, b) \in A$, R is reflexive.	
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Exercises:	
 In 19–31, (1) prove that the relation is an equivalence relation, and (2) describe the distinct equivalence classes of each relation. 19. A is the set of all students at your college. a. R is the relation defined on A as follows: For every x and y in A, 	
$x R y \iff x$ has the same major (or double major) as y.	
 (Assume "undeclared" is a major.) b. S is the relation defined on A as follows: For every x, y ∈ A, 	
$x S y \iff x$ is the same age as y.	
20. E is the relation defined on Z as follows:	
For every $m, n \in \mathbb{Z}$, $m E n \iff 4 (m - n)$.	
H 21. <i>R</i> is the relation defined on \mathbf{Z} as follows:	
For every $m, n \in \mathbb{Z}$, $m R n \iff 7m - 5n$ is even.	103

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SOLUTION	
(a)	
$\begin{aligned} A &= \text{set of all students at your college} \\ R &= \{(x,y) \in A \times A x \text{ has the same major as } y \} \end{aligned}$	
(1) To proof: R is an equivalence relation	Equivalence relation
PROOF	Since R is reflexive, symmetric and transitive, R is also an equivalence relation.
Reflexive	
Since an individual always has the same major as itself, $(x, x) \in R$ if $x \in A$ and thus R is indeed reflexive.	for all
Symmetric	
Let us assume that $(x, y) \in R$.	(3) Each aminglang along will gest in students who have the same main
By the definition of R : x and y have the same major.	(2) Each equivalence class will contain students who have the same major or students who have the same double major.
But then y and x also have the same major:	Thus there is an equivalence class per major and per double major.
$(y,x)\in R$	
Since $(x, y) \in R$ implies $(y, x) \in R$, R is symmetric.	
Transitive	
Let us assume that $(x, y) \in R$ and $(y, z) \in R$.	
By the definition of R :	
x and y have the same major	
y and z have the same major	
But then x and z also have the same major:	
$(x,z)\in R$	
Since $(x, y) \in R$ and $(y, z) \in R$ implies $(x, z) \in R$, R is transitive.	

(b)	
$\begin{split} A &= \text{set of all students at your college} \\ R &= \{(x,y) \in A \times A x \text{ has the same age as } y\} \end{split}$	
(1) To proof: ${\cal R}$ is an equivalence relation	
PROOF	
Reflexive	
Since an individual always has the same age as itself, $(x,x)\in R$ for a and thus R is indeed reflexive.	$\operatorname{il} x \in A$
Symmetric	
Let us assume that $(x, y) \in R$.	Equivalence relation
By the definition of R : x and y have the same age.	Since R is reflexive, symmetric and transitive, R is also an equivalence
But then y and x also have the same age:	relation.
$(y,x)\in R$	
Since $(x, y) \in R$ implies $(y, x) \in R$, R is symmetric.	(2) Each equivalence class will contain students who have the same age.
Transitive	Thus there is an equivalence class per age.
Let us assume that $(x, y) \in R$ and $(y, z) \in R$.	
By the definition of R :	
x and y have the same age	
y and z have the same age	
But then x and z also have the same age:	
$(x,z)\in R$	
Since $(x, y) \in K$ and $(y, z) \in K$ implies $(x, z) \in R$, R is transitive.	

		
23.	Let <i>P</i> be a set of parts shipped to a company from various suppliers. <i>S</i> is the relation defined on <i>P</i> as follows: For every $x, y \in P$,	
	$x S y \iff x$ has the same part number and is shipped from the same supplier as y.	
24.	Let <i>A</i> be the set of identifiers in a computer pro- gram. It is common for identifiers to be used for only a short part of the execution time of a program and not to be used again to execute other parts of the program. In such cases, arranging for identi- fiers to share memory locations makes efficient use of a computer's memory capacity. Define a relation <i>R</i> on <i>A</i> as follows: For all identifiers <i>x</i> and <i>y</i> ,	
	$x R y \iff$ the values of x and y are stored in the same memory location during execution of the program.	106

SOLUTION	
A = Set of identifiers in a computer program	
$R = \{(x,y) \in A \times A x \text{ and } y \text{ are stored in the same memory location} \}$	
(1) To proof: R is an equivalence relation	
PROOF	
Reflexive	
Let $x \in A$	
Since identifier x is always stored in the same memory location as itself, $(x,x)\in R.$	
Since $(x, x) \in R$ for all $x \in A$, R is indeed reflexive.	Equivalence relation
Symmetric	
Let us assume that $(x, y) \in R$. By the definition of R :	Since R is reflexive, symmetric and transitive, R is also an equivalence relation.
x and y are stored in the same memory location	
However, if x and y are stored in the same memory location, then y and x are also stored in the same memory location.	
$(y,x) \in R$	
Since $(x, y) \in R$ implies $(y, x) \in R$, R is symmetric.	(2) Two identifiers are related if they are stored in the same memory location during the execution of the program.
Transitive	
Let us assume that $(x, y) \in R$ and $(y, z) \in R$.	Each equivalence class then contains all identifiers that are stored in the same memory location during the execution of the program
By the definition of R :	come memory receiption during the execution of the program.
\boldsymbol{x} and \boldsymbol{y} are stored in the same memory location	
y and z are stored in the same memory location	
However, this then implies that x , y and z are all stored in the same memory location, thus x and z are stored in the same memory location.	
$(x,z) \in R$	
Since $(x, y) \in R$ and $(y, z) \in R$ implies $(x, z) \in R$, R is transitive.	

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	$x R y \iff$ the values of x and y are stored in the same memory location during execution of the program.	
H	26. <i>D</i> is the relation defined on Z as follows: For every $m, n \in \mathbf{Z}$,	
	$m D n \iff 3 \mid (m^2 - n^2).$	108



