

Each Question 12 points
Exam out of 72 points

22/11/2018

Comp439

[1] (a) What are the programming languages views? Explain.

- 1- Designer (Inventor of the PL)
- 2- Implementor (Who develops & writes the compiler)
- 3- User (who writes programs in this PL)

(b) Explain Briefly the meaning and give examples of **Orthogonality** in a PL.

The PL should behave same thing in similar contexts.

EX: In IBM machines, Addition is performed:

A register, memory
AR register1, register2 > Less orthogonal

But in VAX machines, there is only one instruction.

ADDL operand1, operand2 > more orthogonal
better

(c) What are the programming languages paradigms. Give example on each.

- 1- Imperative (procedural) paradigm: Pascal, C
- 2- Functional Paradigm: Lisp
- 3- Logical Paradigm: Prolog
- 4- Object Oriented Paradigm: Java

12 [2] (a) Write a function in Clisp, **min**(x y) which computes the minimum value of x and y.

```
> (defun min (x y)
    (if (> x y) y x))
```

6

(b) What is the output of the following function in Clisp, justify your answer by **tracing** the function call (func 17 3).

```
> (defun func (a b)
    (if (> a b) 0
        (+ 1 (func (- a b) b))))
```

> (func 17 3) = 5

↓
(1 + (func 14 3)) = 5

↓
(1 + (func 11 3)) = 4

↓
(1 + (func 8 3)) = 3

↓
(1 + (func 5 3)) = 2

↓
(1 + (func 2 3)) = 1

↓
0

6

what the function **func** do?

2

function **func**(m, n) computes & returns the value of $m \text{ div } n$ (m/n)

12 [3] Given the following simple grammar:

block \rightarrow **begin** decls stmts **end**
decls \rightarrow **var** dec-item | λ
dec-item \rightarrow **D** | **D**, dec-item
stmts \rightarrow statement ; stmts | λ
statement \rightarrow **S** | λ

$V_T = \{\text{begin, end, var, D, S, ;, ,}\}$

$V_N = \{\text{block, decls, decl-item, stmts, statement}\}$

(a) Give a program generated by this grammar.

4
begin
var D, D
S;
S;
S;
end

(b) Rewrite production rules using EBNF notations.

4
block \rightarrow begin decls stmts end
decls \rightarrow var dec-item | λ
dec-item \rightarrow D (, D)^{*}
stmts \rightarrow (statement ;)^{*}
statement \rightarrow S | λ

(c) Compute FOLLOW(dec-item).

4
FOLLOW(dec-item) = FOLLOW(decls)
= FIRST(stmts) \cup {end}
= {S, ;} \cup {end} = {S, ;, end}

12 [4] (a) Given the following in C code:

```
int power(int m, int n); // The function power computes and returns m^n
```

```
void main()
{ const int max=10;
  float x=1, y=10;
  x += y;
  int p = power (max,2);
```

Draw the symbol table after the above code is executed.

Name	Type	Value
Power	function Name	100
max	integer Constant	10
x	float Variable	≠ 11
y	float Variable	10
p	integer Variable	100

(b) Given the grammar:

- $G \rightarrow S\$$
- $S \rightarrow AS$
- $A \rightarrow AAB \mid a \mid \lambda$
- $B \rightarrow bBS \mid c \mid \lambda$

	FIRST	FOLLOW
G	a b c λ	—
S	a b c	λ a b c
A	a b c λ	a b c
B	b c λ	a b c

[5] (a) Given the grammar:

$E \rightarrow E + T \mid T$

$T \rightarrow T * F \mid F$

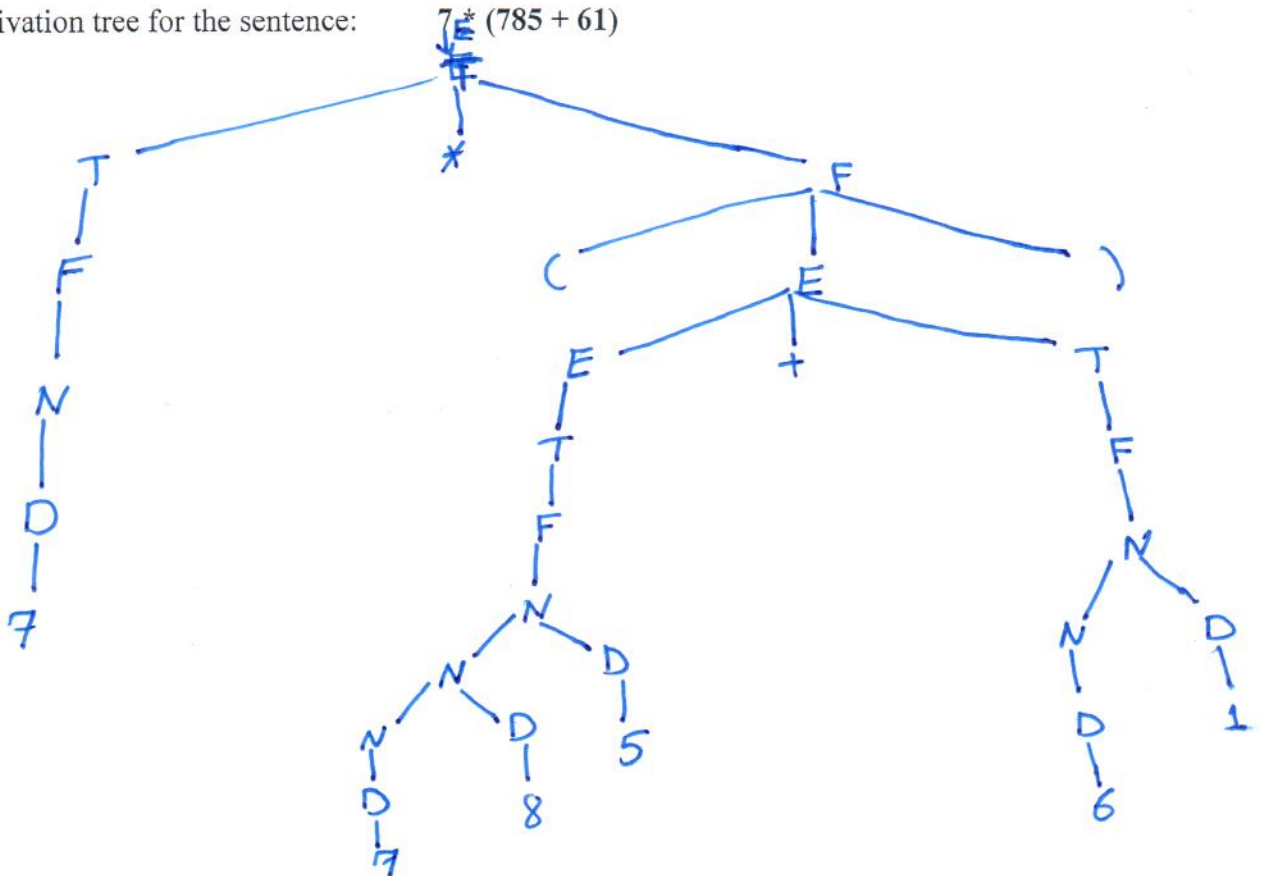
$F \rightarrow (E) \mid N$

$N \rightarrow ND$

$N \rightarrow D$

$D \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$

Draw the derivation tree for the sentence: $7^* (785 + 61)$

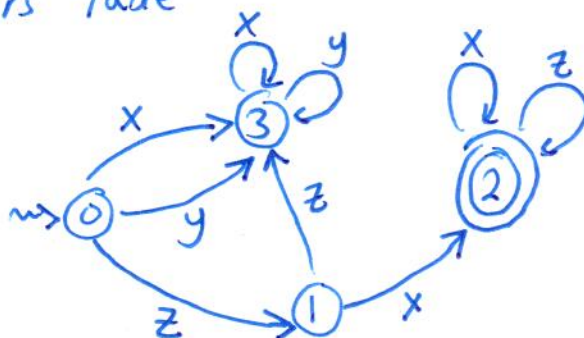


(b) Given the following DFSA. Reduce it to **minimum** states.

	x	y	z
(3,4)	(3,4)	(4,3)	(5,6)
(2,5)	(6,6)		(5,5)
(2,6)	(6,2)		(6,5)
(5,6)	(6,2)		(6,5)

Feasible Pairs Table

$3 \equiv 4$
 $2 \equiv 5 \equiv 6$



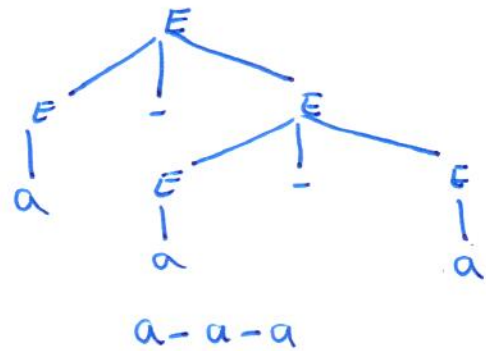
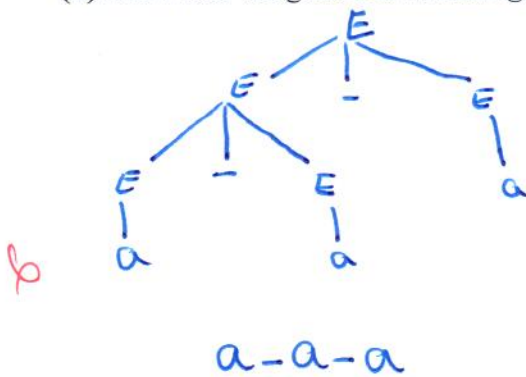
δ	x	y	z
0	3	3	1
1	5		4
2	6		5
3	3	4	
4	4	3	
5	6		6
6	2		5

	x	y	z
0	3	3	1
1	2		3
2	2		2
3	3	3	

6 [6] Given the grammar:

$$E \rightarrow E - E \mid (E) \mid a$$

(a) Show that the grammar is **ambiguous**.



Two derivation trees

(b) A student transform the above grammar to the following none ambiguous grammar:

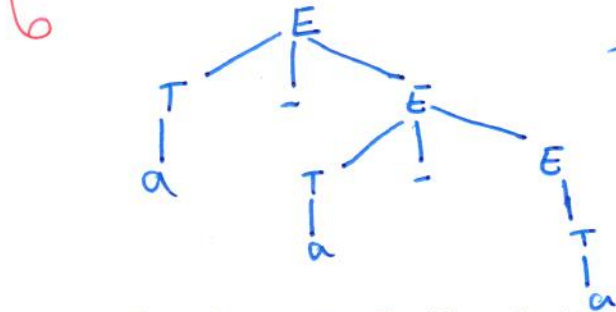
$$E \rightarrow T - E \mid T$$

$$T \rightarrow (E) \mid a$$

What is wrong with code. Explain.

This grammar is right associative which contradicts the associative rules in the "-" operations

consider the sentence a-a-a, its derivation tree is:



This mean that a-a-a will be executed as a-(a-a) which contradict the associativity rule.

(c) Given the grammar G with productions:

$$A \rightarrow \alpha$$

$$A \rightarrow \beta$$

$$A \rightarrow \lambda, \text{ where } \alpha, \beta \neq \lambda$$

State explicitly the conditions so that the above grammar is LL(1).

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