



## FACULTY OF ENGINEERING AND TECHNOLOGY

## COMPUTER SCIENCE DEPARTMENT

COMP233 Discrete Mathematics

# CHAPTER 8

Relations

# Outline

- Relations on Sets
- Properties of Relations
- Equivalence Relations

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- Relations on Sets
- Properties of Relations
- Equivalence Relations

- Partition Sets
- Equivalence Relations
  - Equivalence Relation on a Set of Subsets
  - Equivalence Relation of Digital Logic Circuits
  - Equivalence Relation on a Set of Identifiers
- Equivalence Classes
  - Congruence Modulo n

## Partition Sets

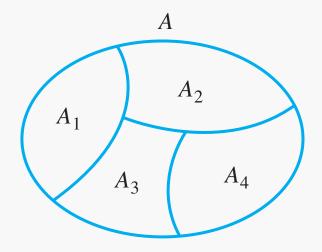
- Equivalence Relations
  - Equivalence Relation on a Set of Subsets
  - Equivalence Relation of Digital Logic Circuits
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  - Relational Numbers

## Partitions of Sets

- A partition of a set is a grouping of its elements into non-empty subsets, in such a way that every element is included in exactly one subset.
- For example, if we divide all the students in this class to groups, where no student is in two groups at the same time, then these groups will be a partition of the class.

# Partition of Sets – Formal Definition

- A finite or infinite collection of non-empty sets  $\{A_1, A_2, A_3, ...\}$  is a partition of a set A if, and only if,
- 1. The set *A* is the union of all the  $A_i s$
- 2. The sets  $A_1, A_2, A_3, \dots$  are mutually disjoint



## Partition of Sets - Example

Let  $\mathbb{Z}$  be the set of all integers, and let

 $T_0 = \{n \in \mathbb{Z} \mid n = 3k, \text{ for some integer } k\},\$  $T_1 = \{n \in \mathbb{Z} \mid n = 3k + 1, \text{ for some integer } k\},\$  $T_2 = \{n \in \mathbb{Z} \mid n = 3k + 2, \text{ for some integer } k\}.$ 

Is  $\{T_0, T_1, T_2\}$  a partition of  $\mathbb{Z}$ ?

Yes. By the quotient-remainder theorem, every integer n can be represented in exactly one of the three forms

n=3k or n=3k+1 or n=3k+2

It also implies that every integer is in one of the sets  $T_0$ ,  $T_1$ ,  $T_2$ , so  $\mathbb{Z} = T_0 \cup T_1 \cup T_2$ .

# **Relations Induced by Partitions**

- A relation induced by a partition is a relation between two elements in the same partition.
- Given a partition of a set *A*, the relation induced by the partition, *R*, is defined on *A* as follows:
- For all  $x, y \in A$ ,

 $x R y \leftrightarrow$  there is a subset  $A_i$  of the partition such that both x and y are in  $A_i$ .

# **Relations Induced by Partitions – Example**

- Let  $A = \{0, 1, 2, 3, 4\}$  and consider the following partition of A:  $\{0, 3, 4\}, \{1\}, \{2\}$
- Find the relation *R* induced by this partition.

Since  $\{0, 3, 4\}$  is a subset of the partition,

0 R 3 because both 0 and 3 are in  $\{0, 3, 4\}$ ,

- 3 R 0 because both 3 and 0 are in  $\{0, 3, 4\}$ ,
- 0 R 4 because both 0 and 4 are in  $\{0, 3, 4\}$ ,
- 4 R 0 because both 4 and 0 are in  $\{0, 3, 4\}$ ,
- 3 R 4 because both 3 and 4 are in  $\{0, 3, 4\}$ ,
- 4 R 3 because both 4 and 3 are in  $\{0, 3, 4\}$ .

# **Relations Induced by Partitions – Example**

- Let  $A = \{0, 1, 2, 3, 4\}$  and consider the following partition of A:  $\{0, 3, 4\}, \{1\}, \{2\}$
- Find the relation *R* induced by this partition.

Also,

- 0 R 0 because both 0 and 0 are in  $\{0, 3, 4\}$
- 3 R 3 because both 3 and 3 are in  $\{0, 3, 4\}$ ,
- 4 *R* 4 because both 4 and 4 are in {0, 3, 4}.

# **Relations Induced by Partitions – Example**

- Let  $A = \{0, 1, 2, 3, 4\}$  and consider the following partition of A:  $\{0, 3, 4\}, \{1\}, \{2\}$
- Find the relation *R* induced by this partition.

Since  $\{1\}$  is a subset of the partition,

1 R 1 because both 1 and 1 are in  $\{1\}$ ,

and since  $\{2\}$  is a subset of the partition,

2 R 2 because both 2 and 2 are in {2}.

Hence

 $R = \{(0,0), (0,3), (0,4), (1,1), (2,2), (3,0), (3,3), (3,4), (4,0), (4,3), (4,4)\}$ 

# Properties of Relations Induced by Partitions

■ Let *A* be a set with a partition and let *R* be the relation induced by the partition. Then *R* is *reflexive*, *symmetric*, and *transitive*.

Proof that *R* is reflexive:

- Suppose  $x \in A$ . Since  $\{A_1, A_2, ..., A_n\}$  is a partition of A, it follows that  $x \in A_i$  for some i.
- But then the statement

there is a set  $A_i$  of the partition such that  $x \in A_i$  and  $x \in A_i$  is true.

Thus, by definition of R, x R x.

# Properties of Relations Induced by Partitions

■ Let *A* be a set with a partition and let *R* be the relation induced by the partition. Then *R* is *reflexive*, *symmetric*, and *transitive*.

Proof that *R* is symmetric:

- Suppose x and y are elements of A such that x R y.
- Then there is a subset  $A_i$  of the partition such that  $x \in A_i$  and  $y \in A_i$  by definition of R.
- It follows that

there is a subset  $A_i$  of the partition such that  $y \in A_i$  and  $x \in A_i$ 

• Hence, by definition of R, y R x.

# **Properties of Relations Induced by Partitions**

■ Let *A* be a set with a partition and let *R* be the relation induced by the partition. Then *R* is *reflexive*, *symmetric*, and *transitive*.

Proof that *R* is transitive:

- Suppose x, y, and z are in A and x R y and y R z. By definition of R, there are subsets  $A_i$  and  $A_j$  of the partition such that x and y are in  $A_i$  and y and z are in  $A_j$ .
- Suppose  $A_i \neq A_j$ . Then  $A_i \cap A_j = \emptyset$  since  $\{A_1, A_2, \dots, A_n\}$  is a partition of A.
- But y is in  $A_i$  and y is also in  $A_j$ . Hence  $A_i \cap A_j \neq \emptyset$ . Thus,  $A_i = A_j$  (by contradiction).
- It follows that x, y, and z are all in  $A_i$ , and so in particular, x and z are in  $A_i$ .
- Thus, by definition of R, x R z.

Partition Sets

### Equivalence Relations

- Equivalence Relation on a Set of Subsets
- Equivalence Relation of Digital Logic Circuits
- Equivalence Relation on a Set of Identifiers
- Equivalence Classes
  - Congruence Modulo n
  - Relational Numbers

A relation on a set that satisfies the three properties of reflexivity, symmetry, and transitivity is called an equivalence relation.

#### Definition:

Let *A* be a set and *R* a relation on *A*. *R* is an equivalence relation if, and only if, *R* is reflexive, symmetric, and transitive.

Thus, according to the theorem that we have just proved, the relation induced by a partition is an equivalence relation.

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- Let X be the set of all nonempty subsets of  $\{1, 2, 3\}$ . Then  $X = \{\{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$
- Define a relation R on X as follows: For all A and B in X,  $A R B \iff the \ least \ element \ of A \ equals \ the \ least \ element \ of B$
- Prove that R is an equivalence relation on X.
- *R* is reflexive:

Suppose *A* is a nonempty subset of  $\{1, 2, 3\}$ . We must show that *A R A*. It is true to say that the least element of *A* equals the least element of *A*. Thus, by definition of *R*, *A R A*.

- Let X be the set of all nonempty subsets of  $\{1, 2, 3\}$ . Then  $X = \{\{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$
- Define a relation R on X as follows: For all A and B in X,  $A R B \iff the \ least \ element \ of A \ equals \ the \ least \ element \ of B$
- Prove that R is an equivalence relation on X.
- *R* is symmetric:

Suppose A and B are nonempty subsets of {1, 2, 3} and A R B. We must show that B R A.

Since *A R B*, the least element of *A* equals the least element of *B*. But this implies that the least element of *B* equals the least element of *A*, and so, by definition of *R*, *B R A*.

- Let X be the set of all nonempty subsets of  $\{1, 2, 3\}$ . Then  $X = \{\{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$
- Define a relation R on X as follows: For all A and B in X,  $A R B \Leftrightarrow the least element of A equals the least element of B$
- Prove that R is an equivalence relation on X.
- *R* is transitive:

Suppose A, B, and C are nonempty subsets of {1, 2, 3}, A R B, and B R C. We must show that A R C.

Since *A R B*, the least element of *A* equals the least element of *B*, and since *B R C*, the least element of *B* equals the least element of *C*.

Thus, the least element of A equals the least element of C, and so, by definition of R, A R C.

- Let X be the set of all nonempty subsets of  $\{1, 2, 3\}$ . Then  $X = \{\{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$
- Define a relation R on X as follows: For all A and B in X,  $A R B \iff the \ least \ element \ of A \ equals \ the \ least \ element \ of B$
- Prove that R is an equivalence relation on X.
- Since R is reflexive, symmetric, and transitive, then R is an equivalence relation on X.

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■ Let *S* be the set of all digital logic circuits with a fixed number *n* of inputs. Define a relation *E* on *S* as follows:

For all circuits  $C_1$  and  $C_2$  in S,

 $C_1 E C_2 \leftrightarrow C_1$  has the same input/output table as  $C_2$ .

- If  $C_1 E C_2$ , then circuit  $C_1$  is said to be equivalent to circuit  $C_2$ .
- Prove that E is an equivalence relation on S.

For all circuits  $C_1$  and  $C_2$  in S,

 $C_1 E C_2 \leftrightarrow C_1$  has the same input/output table as  $C_2$ .

■ *E* is reflexive:

Suppose *C* is a digital logic circuit in *S*. We must show that *C E C*.

Certainly, *C* has the same input/output table as itself.

Thus, by definition of *E*, *C E C* as was to be shown.

For all circuits  $C_1$  and  $C_2$  in S,

 $C_1 E C_2 \leftrightarrow C_1$  has the same input/output table as  $C_2$ .

■ *E* is symmetric:

Suppose  $C_1$  and  $C_2$  are digital logic circuits in S such that  $C_1 E C_2$ . We must show that  $C_2 E C_1$ .

By definition of *E*, since  $C_1 E C_2$ , then  $C_1$  has the same input/output table as  $C_2$ .

It follows that  $C_2$  has the same input/output table as  $C_1$ .

Hence, by definition of *E*,  $C_2 E C_1$  as was to be shown.

For all circuits  $C_1$  and  $C_2$  in S,

 $C_1 E C_2 \leftrightarrow C_1$  has the same input/output table as  $C_2$ .

#### ■ *E* is transitive:

Suppose  $C_1$ ,  $C_2$ , and  $C_3$  are digital logic circuits in S such that  $C_1 E C_2$  and  $C_2 E C_3$ . We must show that  $C_1 E C_3$ .

By definition of *E*, since  $C_1 E C_2$  and  $C_2 E C_3$ , then

 $C_1$  has the same input/output table as  $C_2$ 

and  $C_2$  has the same input/output table as  $C_3$ .

It follows that  $C_1$  has the same input/output table as  $C_3$ 

Hence, by definition of *E*,  $C_1 E C_3$  as was to be shown.

For all circuits  $C_1$  and  $C_2$  in S,

 $C_1 E C_2 \leftrightarrow C_1$  has the same input/output table as  $C_2$ .

Since *E* is reflexive, symmetric, and transitive, then *E* is an equivalence relation on *S*.

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■ Let *L* be the set of all allowable identifiers in a certain computer language, and define a relation *R* on *L* as follows: For all strings *s* and *t* in *L*,

 $s \ R \ t \leftrightarrow$  the first eight characters of s equal the first eight characters of t.

• Prove that R is an equivalence relation on L.

*R* is reflexive:

- Let  $s \in L$ . We must show that s R s.
- Clearly s has the same first eight characters as itself. Thus, by definition of R, s R s as was to be shown.

*R* is symmetric:

- Let s and t be in L and suppose that s R t. We must show that t R s.
- By definition of R, since s R t, the first eight characters of s equal the first eight characters of t.
- But then the first eight characters of t equal the first eight characters of s.
- And so, by definition of R, t R s as was to be shown.

*R* is transitive:

- Let *s*, *t*, and *u* be in *L* and suppose that *s R t* and *t R u*. We must show that *s R u*.
- By definition of R, since s R t and t R u,

the first eight characters of s equal the first eight characters of t,

and the first eight characters of t equal the first eight characters of u.

- Hence the first eight characters of s equal the first eight characters of u.
- Thus, by definition of R, s R u as was to be shown.

■ Let *L* be the set of all allowable identifiers in a certain computer language, and define a relation *R* on *L* as follows: For all strings *s* and *t* in *L*,

 $s \ R \ t \leftrightarrow$  the first eight characters of s equal the first eight characters of t.

- Prove that R is an equivalence relation on L.
- Since R is reflexive, symmetric, and transitive, then R is an equivalence relation on L.

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# Equivalence Classes of an Equivalence Relation

- Suppose A is a set and R is an equivalence relation on A. For each element a in A, the equivalence class of a, denoted [a] and called the class of a for short, is the set of all elements x in A such that x is related to a by R.
- In symbols:

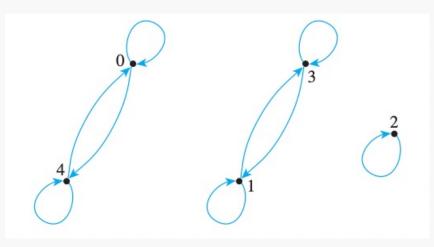
$$[a] = \{x \in A \mid x \mathrel{R} a\}$$

- When several equivalence relations on a set are under discussion, the notation  $[a]_R$  is often used to denote the equivalence class of a under R.
- The procedural version of this definition is

for all 
$$x \in A, x \in [a] \leftrightarrow x R a$$

#### Equivalence Classes of a Relation – Example

- Let  $A = \{0, 1, 2, 3, 4\}$  and define a relation R on A as follows:  $R = \{(0, 0), (0, 4), (1, 1), (1, 3), (2, 2), (3, 1), (3, 3), (4, 0), (4, 4)\}$
- Draw the directed graph of R.
- Does this graph show that *R* is an equivalence relation?
   Yes, it does.



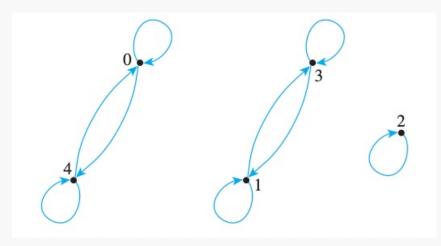
#### Equivalence Classes of a Relation – Example

Let 
$$A = \{0, 1, 2, 3, 4\}$$
 and define a relation  $R$  on  $A$  as follows:  
 $R = \{(0, 0), (0, 4), (1, 1), (1, 3), (2, 2), (3, 1), (3, 3), (4, 0), (4, 4)\}$ 

Find the distinct equivalence classes of R.

First find the equivalence class of every element of *A*.

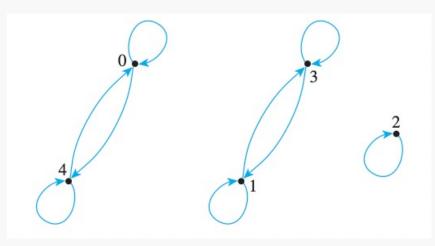
 $[0] = \{x \in A \mid x R \ 0\} = \{0, 4\}$   $[1] = \{x \in A \mid x R \ 1\} = \{1, 3\}$   $[2] = \{x \in A \mid x R \ 2\} = \{2\}$   $[3] = \{x \in A \mid x R \ 3\} = \{1, 3\}$  $[4] = \{x \in A \mid x R \ 4\} = \{0, 4\}$ 



#### Equivalence Classes of a Relation – Example

• Let  $A = \{0, 1, 2, 3, 4\}$  and define a relation R on A as follows:  $R = \{(0, 0), (0, 4), (1, 1), (1, 3), (2, 2), (3, 1), (3, 3), (4, 0), (4, 4)\}$ 

Find the distinct equivalence classes of *R*.
 Note that [0] = [4] and [1] = [3].
 Thus, the distinct equivalence classes of the relation are {0, 4}, {1, 3}, and {2}.



## Lemmas of Equivalence Classes

- Suppose *A* is a set, *R* is an equivalence relation on *A*, and *a* and *b* are elements of *A*. If a R b, then [a] = [b].
- If *A* is a set, *R* is an equivalence relation on *A*, and *a* and *b* are elements of *A*, then either  $[a] \cap [b] = \emptyset$  or [a] = [b].
- If A is a set and R is an equivalence relation on A, then the distinct equivalence classes of R form a partition of A; that is, the union of the equivalence classes is all of A, and the intersection of any two distinct classes is empty.

### **Class Representative**

- Suppose *R* is an equivalence relation on a set *A* and *S* is an equivalence class of *R*.
- A representative of the class S is any element a such that [a] = S.

# **Equivalence Relations**

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- Let *R* be the relation of congruence modulo 3 on the set  $\mathbb{Z}$  of all integers.
- That is, for all integers *m* and *n*,  $m R n \Leftrightarrow 3 \mid (m - n) \leftrightarrow m \equiv n \pmod{3}$
- Describe the distinct equivalence classes of R.
- For each integer a,  $[a] = \{x \in Z | x R a\}$   $= \{x \in Z | 3| (x - a)\}$   $= \{x \in Z | x - a = 3k, for some integer k\}$

■ Therefore,

 $[a] = \{x \in Z \mid x = 3k + a, for some integer k\}$ 

In particular,

$$[0] = \{x \in Z \mid x = 3k + 0, for some integer k\} \\ = \{x \in Z \mid x = 3k, for some integer k\} \\ = \{... - 9, -6, -3, 0, 3, 6, 9, ... \}$$

 $[1] = \{x \in Z \mid x = 3k + 1, for some integer k\} \\ = \{\dots - 8, -5, -2, 1, 4, 7, 10, \dots\}$ 

 $[2] = \{x \in Z \mid x = 3k + 2, for some integer k\} \\= \{\dots - 7, -4, -1, 2, 5, 8, 11, \dots\}$ 

• Now since 3 R 0, then [3] = [0].

- More generally, by the same reasoning,  $[0] = [3] = [-3] = [6] = [-6] = \cdots$ , and so on.
- Similarly,  $[1] = [4] = [-2] = [7] = [-5] = \cdots$ , and so on.
- And  $[2] = [5] = [-1] = [8] = [-4] = \cdots$ , and so on.
- Notice that every integer is in class [0], [1], or [2]. Hence the distinct equivalence classes are

 $\{x \in Z \mid x = 3k, for some integer k\},\$  $\{x \in Z \mid x = 3k + 1, for some integer k\}, and\$  $\{x \in Z \mid x = 3k + 2, for some integer k\}.$ 

the distinct equivalence classes are

 $\{x \in Z \mid x = 3k, for some integer k\},\$  $\{x \in Z \mid x = 3k + 1, for some integer k\}, and\$  $\{x \in Z \mid x = 3k + 2, for some integer k\}.$ 

- In words, the three classes of congruence modulo 3 are
- (1) the set of all integers that are divisible by 3,
- (2) the set of all integers that leave a remainder of 1 when divided by 3, and
- (3) the set of all integers that leave a remainder of 2 when divided by 3.

# Congruence Modulo *n*

- Let m and n be integers and let d be a positive integer.
- We say that m is congruent to n modulo d and write

 $m \equiv n \pmod{d}$  if, and only if,  $d \mid (m - n)$ .

Symbolically:

$$m \equiv n \pmod{d} \leftrightarrow d \mid (m - n)$$

# **Evaluating Congruences**

- Determine which of the following congruences are true and which are false.
- a.  $12 \equiv 7 \pmod{5}$ 
  - True.  $12 7 = 5 = 5 \cdot 1$ . Hence  $5 \mid (12 7)$ , and so  $12 \equiv 7 \pmod{5}$ .
- b.  $6 \equiv -8 \pmod{4}$ 
  - False. 6 (-8) = 14, and 4  $\downarrow$ 14 because 14  $\neq$  4  $\cdot$  k for any integer k. Consequently, 6  $\not\equiv$  -8 (mod 4).
- c.  $3 \equiv 3 \pmod{7}$ 
  - True.  $3 3 = 0 = 7 \cdot 0$ . Hence  $7 \mid (3 3)$ , and so  $3 \equiv 3 \pmod{7}$ .

# **Equivalence Relations**

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#### **Rational Numbers As Equivalence Classes**

 Let A be the set of all ordered pairs of integers for which the second element of the pair is nonzero. Symbolically,

 $A = Z \times (Z - \{0\})$ 

- Define a relation *R* on *A* as follows: For all  $(a, b), (c, d) \in A$ ,  $(a, b)R (c, d) \leftrightarrow ad = bc$
- The fact is that *R* is an equivalence relation.

Describe the distinct equivalence classes of R.

#### **Rational Numbers As Equivalence Classes**

- There is one equivalence class for each distinct rational number.
- Each equivalence class consists of all ordered pairs (a, b) that, if written as fractions  $\frac{a}{b}$ , would equal each other.
- The reason for this is that the condition for two rational numbers to be equal is the same as the condition for two ordered pairs to be related.
  - For instance, the class of (1, 2) is  $[(1,2)] = \{(1,2), (-1,-2), (2,4), (-2,-4), (3,6), (-3,-6), ...\}$ since  $\frac{1}{2} = \frac{-1}{-2} = \frac{2}{4} = \frac{-2}{-4} = \frac{3}{6} = \frac{-3}{-6}$  and so forth.