



## ANSWER BOOKLET

Student: <u>Digital</u> Number <u>3</u>
Course: Department: ..... Number: .....
Division: ..... Instructor: .....
Date: ..... Day Month Year

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### For Instructor's Use

Question	Grade
1	
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<b>Total</b>	

## Examples of Algebraic Manipulation

$$\textcircled{1} x(x'+y) = xx' + xy = 0 + xy = xy$$

$$\textcircled{2} x + x'y = (x + x'y)(x+y) = 1(x+y) = x+y$$

$$\textcircled{3} (x+y)(x+y')$$

$$= x + xy + x'y + yy' = x(1+y+y') = x$$

$$\textcircled{4} xy + x'z + yz = xy + x'z + yz(x+x')$$

$$= xy + x'z + x'yz + xy'z$$

$$= xy(1+z) + x'z(1+y)$$

$$= xy + x'z$$

$$\textcircled{5} (x+y)(x'+z)(y+z) = (x+y)(x'+z)$$

$$= (xx' + xz + x'y + yz)(y+z)$$

$$= xy^2 + xz^2 + x'y + x'y^2 + yz^2$$

$$= yz^2(x+x') + xz^2 + x'y + yz^2$$

$$= yz^2 + xz^2 + x'y$$

$$= yz^2 + x'y$$

$$= (x+y)z^2$$

$$~~(x+y)(y+z)(x+z)(z+z) + x'y~~$$

$$~~x+x'~~$$

$$= yz^2(x+x') + xz^2 + x'y =$$

$$xz^2 + x'y^2 + xz^2 + x'y = xz^2(y+1) + x'y(z+1)$$

$$= xz^2 + x'y = ~~(x+x')(x+y)(x'+z)~~$$

$$= (x+y)(x'+z)$$

## ⊛ Complement of a function

Demorgan's theorems :-

$$(A+B+C+\dots+Z)' = A'B'C'D'\dots Z'$$

$$(ABCDE\dots Z)' = A'+B'+C'+D'+E'+\dots+Z'$$

Example

$$F = x'yz' + x'y'z$$

$$\begin{aligned}\rightarrow F' &= (x'yz' + x'y'z)' \\ &= (x+y+z)(x+y+z')\end{aligned}$$

Ex

$$F = x(x'z' + yz)$$

$$\begin{aligned}F' &= (x(x'z' + yz))' \\ &= x' + (y'z' + yz)' \\ &= x' + (y'z')' \cdot (yz)' \\ &= x' + (y+z)(y'+z')\end{aligned}$$

2-5

## ④ Canonical and Standard Forms

- a binary variable ( $x$ ) may ~~be~~ appear in its normal form ( $x$ ) or in its complement form ( $x'$ ).
- for two binary variables ( $x, y$ ) <sup>combined with AND operation</sup>  $\rightarrow$  there are  $2^2$  forms:  $x'y'$ ,  $x'y$ ,  $xy'$ ,  $xy$ .
- these terms are called minterms or a standard product.
- for  $n$  variables  $\rightarrow 2^n$  forms (minterms).
- for two binary variables ( $x, y$ ), combined with OR operation  $\rightarrow$  there are  $2^n$  forms:  $x+y$ ,  $x+y'$ ,  $x'+y$ ,  $x'+y'$ .
- these terms are called maxterms or standard sum.

$x$	$y$	<u>Minterm</u>		<u>Maxterm</u>	
		Term	Designation	Term	Designation
0	0	$x'y'$	$m_0$	$x+y$	$M_0$
0	1	$x'y$	$m_1$	$x+y'$	$M_1$
1	0	$xy'$	$m_2$	$x'+y$	$M_2$
1	1	$xy$	$m_3$	$x'+y'$	$M_3$

⊙ sum of all minterms (for 2 variables  $x, y$ )  
 $\Rightarrow x'y' + x'y + xy' + xy = 1$  (always)

⊙ product of all maxterms (for 2 variables  $x, y$ )  
 $\Rightarrow (x+y) \cdot (x+y') \cdot (x'+y) \cdot (x'+y') = 0$  (always)

Ex.

x	y	z	minterms		maxterms	f <sub>1</sub>	f <sub>2</sub>
			f <sub>1</sub>	f <sub>2</sub>			
0	0	0	$x'y'z'$ (m <sub>0</sub> )	$x+y+z$	0	0	
0	0	1	$x'y'z$ (m <sub>1</sub> )	$x+y+z'$	1	0	
0	1	0	$x'y z'$ (m <sub>2</sub> )	$x'+y+z$	0	0	
0	1	1	$x'y z$ (m <sub>3</sub> )	$x'+y+z'$	0	1	
1	0	0	$x y' z'$ (m <sub>4</sub> )	$x'+y+z$	1	0	
1	0	1	$x y' z$ (m <sub>5</sub> )	$x'+y+z'$	0	1	
1	1	0	$x y z'$ (m <sub>6</sub> )	$x'+y'+z$	0	1	
1	1	1	$x y z$ (m <sub>7</sub> )	$x'+y'+z'$	1	1	

express f<sub>1</sub> in term of minterms

$$\Rightarrow f_1 = x'y'z + xy'z' + xyz$$

$$= m_1 + m_4 + m_7 = \sum (1, 4, 7)$$

in the same way f<sub>2</sub> = m<sub>3</sub> + m<sub>5</sub> + m<sub>6</sub> + m<sub>7</sub>

$$= \sum (3, 5, 6, 7)$$

express f<sub>1</sub> in term of maxterms

$$f_1 = (x+y+z) \cdot (x+y'+z) \cdot (x'+y+z')$$

$$= M_0 \cdot M_2 \cdot M_5 \cdot M_6 = \prod (0, 2, 3, 5, 6)$$

$$f_2 = M_0 M_1 M_2 M_4$$

④ Any Boolean function can be expressed as a sum of ~~the~~ minterms or as a product of max terms.

Ex. Express  $f = A + B'c$  in a sum of minterms

$$\begin{aligned} \Rightarrow f &= A(B+B') + B'c(A+A') \\ &= AB + AB' + AB'c + A'B'c \\ &= AB(c+c') + AB'(c+c') + AB'c + A'B'c \\ &= \underline{ABC} + ABC' + AB'c + AB'c' + AB'c + A'B'c \\ &= m_7 + m_6 + m_5 + m_4 + m_1 \end{aligned}$$

$$\rightarrow F(A, B, c) = \sum (1, 4, 5, 6, 7)$$

Ex. Express  $f = xy + x'z$  in a product of maxterms  
— change it into OR terms using distributive law

$$\Rightarrow f = xy + x'z$$

$$= (xy + x') \cdot (xy + z)$$

$$= (x + x') \cdot (y + x') \cdot (x + z) \cdot (y + z)$$

$$= (x' + y) \cdot (x + z) \cdot (y + z)$$

$$= (x' + y + zz') \cdot (x + z + yy') \cdot (y + z + xx')$$

$$= (x' + y + z) \cdot (x' + y + z') \cdot (x + z + y) \cdot (x + z + y')$$

$$(\cancel{y + z + x}) \cdot (\cancel{y + z + x'})$$

$$= (x' + y + z) \cdot (x' + y + z') \cdot (x + y + z) \cdot (x + y' + z)$$

$$= M_4 \cdot M_5 \cdot M_0 \cdot M_2 = \Pi(0, 2, 4, 5)$$

## Conversion between Canonical forms

complement of a function:

e.g.  $F(A, B, C) = \sum(1, 4, 5, 6, 7)$

$$\Rightarrow F'(A, B, C) = \sum(0, 2, 3) = m_0 + m_2 + m_3$$

The complement of a function expressed as the sum of minterms equals the sum of minterms missing from the original function.

Use De Morgan's theorem  $\Rightarrow$

$$\begin{aligned} F'(A, B, C) &= (m_0 + m_2 + m_3)' \\ &= m_0' \cdot m_2' \cdot m_3' = M_0 \cdot M_2 \cdot M_3 \\ &= \prod(0, 2, 3) \end{aligned}$$

$m_j' = M_j$

- Using truth table instead of algebraic theorems ( $F = xy + x'z$ )

x	y	z	x'	xy	x'z	xy + x'z	
0	0	0	1	0	0	0	
0	0	1	1	0	1	1	$\Rightarrow F = m_1 + m_3 + m_6 + m_7$
0	1	0	1	0	0	0	$= \sum(1, 3, 6, 7)$
0	1	1	1	0	1	1	$\Rightarrow$
1	0	0	0	0	0	0	$F = \prod(0, 2, 4, 5)$
1	0	1	0	0	0	0	
1	1	0	0	1	0	1	(Same result as using algebra)
1	1	1	0	1	0	1	

## ⊛ Standard forms

- two types of standard forms: sum of products and products of sums
- sum of products

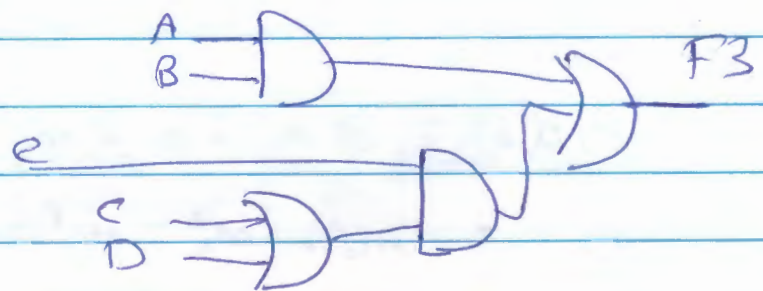
$$F1 = y' + xy + x'y'z'$$

- products of sums

$$F2 = x(y'+z)(x'+y+z')$$

- nonstandard form

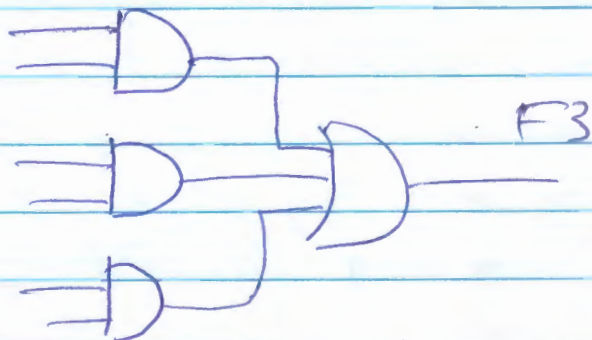
$$F3 = AB + C(D+E)$$



3-level implementation

→ change nonstandard form to a standard form

$$\Rightarrow F3 = AB + CD + CE$$



2-level implementation  
(better).




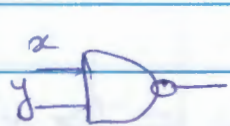
## Digital Logic Gates

① AND   $F$

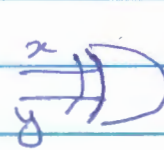
② OR 

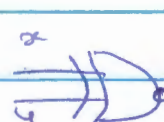
③ Inverter 

④ Buffer 

⑤ NAND   $F = (xy)'$   $\uparrow$

⑥ NOR   $F = (x+y)'$   $\downarrow$

⑦ Exclusive-OR  
XOR   $F = xy' + x'y$   
 $= x \oplus y$

⑧ Exclusive-Nor  
or  
equivalence   $F = xy + x'y'$   
 $= (x \oplus y)'$

x	y	F
0	0	1
0	1	0
1	0	0
1	1	1

## notes

$$(x \downarrow y) \downarrow z \neq x \downarrow (y \downarrow z)$$

$$[(x \downarrow y)' + z]' \neq [x + (y + z)']'$$

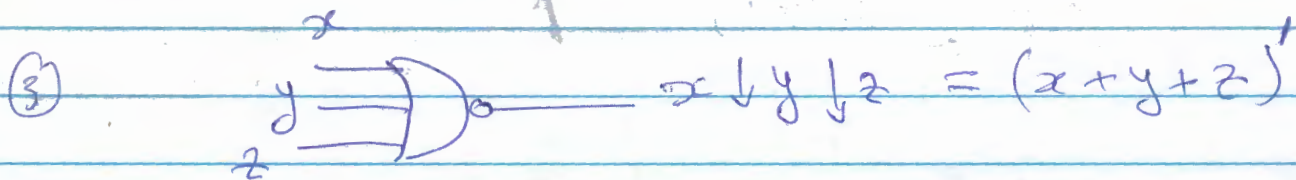
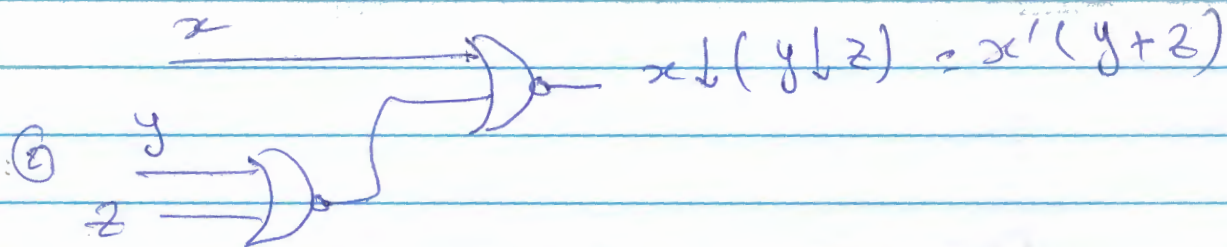
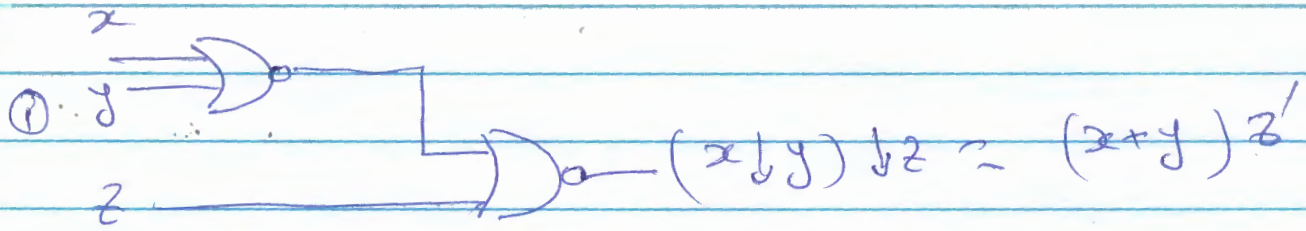
$$xz' + yz' \neq x'y + x'z$$

$\Rightarrow$  By definition

$$x \downarrow y \downarrow z = (x + y + z)'$$

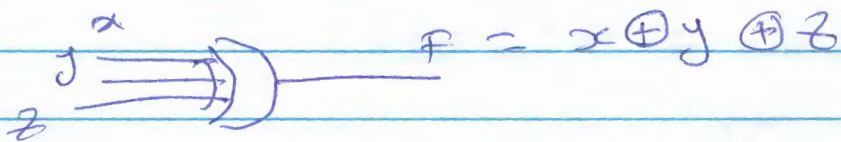
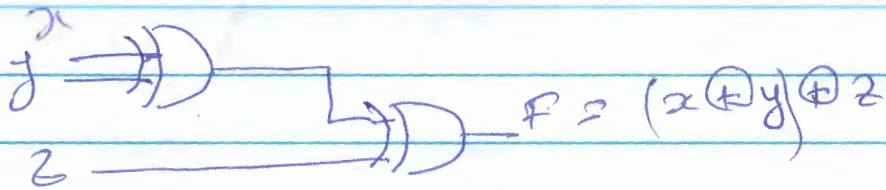
$$x \uparrow y \uparrow z = (xyz)'$$





XOR

$$x \oplus (y \oplus z) = (x \oplus y) \oplus z = x \oplus y \oplus z$$



## CHAPTER 3

⊕ Two variable map

⊕ three variable map

⊕ four variable map

wx \ yz	00	01	11	10
00				
01				
11				
10				

- If we have  $n$  variables
- Two adjacent squares represent a term of all variables except 1
- 4 adjacent squares " " " " variables except 2  $\Rightarrow$  no. of literals =  $n-2$
- $2^m$  adjacent squares " " " " variables except  $m$   $\Rightarrow$  no. of literals =  $n-m$

Ex:  $F(w, x, y, z) = \sum (0, 1, 2, 4, 5, 6, 8, 9, 12, 13, 14)$

		yz				
wx	00	01	11	10		
00	1	1		1	w x y z F	
01	1	1		1	0 0 0 0 1	
11	1	1		1	0 0 0 1 1	
10	1	1		1	0 0 1 0 1	

$\Rightarrow F = y' + w'z' + xz'$

~~Prime~~

		yz				
wx	00	01	11	10		
00	1	1	1	1	$F = y' + w'z' + xz'$	
01	1	1	1	1		
11						
10	1	1	1			

Prime Implicants

- A prime implicant is a product term obtained by combining the maximum possible number of adjacent squares in the map.
- if a minterm in a square is covered by only one prime implicant, that prime implicant is said to be essential.

Ex.

$$F(A, B, C, D) = \sum (0, 2, 3, 5, 7, 8, 9, 10, 11, 13, 15)$$

AB \ CD	00	01	11	10
00	1		1	1
01		1	1	
11		1	1	
10	1	1	1	1

$$\begin{aligned} F &= BD + B'D' + CD + AD \\ &= BD + B'D' + CD + AB' \\ &= BD + B'D' + B'C + AD \\ &= BD + B'D' + B'C + AB' \end{aligned}$$

Ex.  $F(A, B, C) = \sum (0, 1, 2, 5, 7)$

A \ Bc	00	01	11	10
0	1	1		1
1		1	1	

$$\begin{aligned} F &= A'C' + AC + A'B' \\ &= A'C' + AC + B'C \end{aligned}$$

### ⊕ Five Variables Map

AB \ CDE	000	001	011	010	110	111	101	100
00								
01								
11								
10								

Ex:

$$F(A, B, C, D, E) = \sum (0, 2, 4, 6, 9, 13, 21, 23, 25, 29, 31)$$

AB \ CDE	000	001	011	010	110	111	101	100
00	1			1	1			1
01		1					1	
11		1				1	1	
10						1	1	

A = 0					A = 1				
B/C \ D/E	00	01	11	10	B/C \ D/E	00	01	11	10
00	1			1	00	1	1	1	1
01	1			1	01	1	1	1	1
11		1			11	1	1	1	1
10		1			10	1	1	1	1

$$F = ACE + A'B'E' + BD'E$$

### ⊕ Product of Sums Simplification

$$F(A, B, C, D) = \sum (0, 1, 2, 5, 8, 9, 10)$$

<del>AB</del> CD	00	01	11	10
00	1	1	1	1
01	1	1		
11	1	1		
10	1	1		

$$F = B'C' + B'D' + A'C'D$$

$$F' = AB + CD + BD'$$

$$\Rightarrow F = (A' + B')(C' + D')(B' + D)$$

## ⊕ Don't Care Conditions

→ Sometimes, for some combinations of the inputs, the output <sup>has</sup> not specified value, this is called don't care condition and designated by 'X'.

- Don't care conditions help to simplify the function.

Ex.  $F(w, x, y, z) = \sum (1, 3, 7, 11, 15)$

$d(w, x, y, z) = \sum (0, 2, 5)$

$w/x \backslash yz$	00	01	11	10
00	X	1	1	X
01		X	1	
11			1	
10			1	

$F = yz + w'x'$

$F = yz + w'z$

\* note: they don't have an identical truth table!

Ex.  $F(w, x, y, z) = \sum (1, 3, 7)$

$d(w, x, y, z) = \sum (0, 5)$

$w/x \backslash yz$	00	01	11	10
00	X	1	1	
01		X	1	
11				
10				

$F = w'z$